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ELECTROMAGNETIC \& SENSOR SYSTEMS DEPARTMENT
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## Identifying Aliases in Graphs

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May 21, 2008


Quantitative Methods in Defense and National Security

## Outline

Motivation

Definitions and Model

Alias Identification

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## Alias Identification

## Conclusions

## Social Networks

- A model of the relationships between entities.
- Also used to study insurgent groups, terrorist cells, etc.
- Relates actors (nodes in the network) through relationships (edges in the network).
- Typically used for small groups, with full knowledge of all links.



## Covert Networks

- Actors have a vested interest in not being observed.
- Networks may be very large.
- The networks change in time.
- Some links are known to be there, some known to be missing, but others are unknown.
- An actor may try to hide (change email address, change phone number, start calling themselves Colonel Guapa).


## Methodology

- Assume the existence of a "social space" $\mathcal{S}$ which controls the structure of the network.
- The probability of an edge in the network is a function of the "closeness" of the nodes in $\mathcal{S}$.
- The social space provides a framework from which inference can be performed.


## Social Space

- Early work reported by Hoff et al in JASA.
- Model based on location:
- Probability of an edge between $v_{i}$ and $v_{j}$ a function of their distance in social space.
- Several variations proposed.
- Versions of the Exponential Random Graph Models (ERGMs) (Hunter et al, JASA 2008) can be thought of in terms of a "social space".
- We will discuss a "social space" model that has a simple least squares algorithm for fitting the parameters, which can be used on large graphs (thousands to tens of thousands of nodes or more).


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## Graph Definitions

- A graph is a pair $(V, E)$ where $V$ is a set (vertices) and $E$ is a collection of unordered pairs of vertices (edges).
- We can consider directed graphs $(V, A)$ where $A$ (arcs or arrows) are ordered pairs.
- The order of the graph is $|V|$ and the size of the graph is $|E|$ (or $|A|$ in the case of directed graphs (digraphs)).
- Vertices are sometimes called "nodes" or "actors".
- Edges are sometimes called "links" or "relations".
- The adjacency matrix $A=\left(a_{i j}\right)$ is the $|V| \times|V|$ binary matrix with a 1 in those places where an edge occurs in the graph.


## Probabilistic Framework

- We place a probability structure on the network.
- This means we fit a generative model to the graph.
- This allows us to estimate the probability of a missing (unknown) link.
- We can bring node attributes into the model.
- We are essentially choosing the "most likely" graph given the model assumption and the observed edges.


## Random Dot Product Graphs

- Each vertex $v_{i}$ has associated with it a vector $x_{i}$.
- Place an edge $v_{i} v_{j}$ between vertices $v_{i}$ and $v_{j}$ with probability proportional to $x_{i} x_{j}$, the dot product of $x_{i}$ and $x_{j}$.
- Thus $p_{i j}=f\left(x_{i} x_{j}\right)$. We'll use the threshold function for $f$ :

$$
f(x)= \begin{cases}0 & x<0 \\ x & 0 \leq x \leq 1 \\ 1 & x>1\end{cases}
$$

- The edges in the random graph are no longer independent.
- We need to estimate the $x_{i}$ from the observed graph.
- We can extend the model to directed graphs by having inand out-vectors $x_{i}^{\prime}$ and $x_{i}^{O}$ with $p_{i j}$ proportional to $x_{i}^{O} x_{j}^{\prime}$.


## $\mathcal{S}$

- Each vertex $v_{i}$ has associated with it a vector $x_{i} \in \mathcal{S}$.
- The proximity (as measured by the dot product) of two vectors controls the probability of an edge.
- Thus $\mathcal{S}$ is the space which defines the random graph that we observe.
$\mathcal{S}$

$\mathcal{G}$



## Linear Algebra (Least Squares)

Note that if we want to find the vectors $U$ which best "match" the adjacency matrix $A$ (best in Frobenius norm), then the singular value decomposition: $A=U D V^{\prime}$ almost works (the problem is the diagonal). Note that for graphs $A$ is symmetric, so $V=U$.

1. Set $D=\operatorname{diag}(0)$.
$1.1 s=\operatorname{svd}(A+D)$.
$1.2 X=s \$ U$, scaled by the singular values.
$1.3 D=\operatorname{diag}\left(X X^{\prime}\right)$.
2. Repeat 1-3 until convergence.
3. Return $X$.

## The Enron Data

- Graphs (directed graphs) of emails between executives at Enron.
- 184 email addresses (nodes).
- 150 executives (names).
- 187 weeks.
- Each graph corresponds to 1 week of emails.
- An edge $v \rightarrow w$ if there was an email from $v$ to $w$ within the week.
- Note: we are ignoring multiple emails and an email from one to many generates a "star" of edges.


## An Alias



Computer scientists are analyzing about a half million Enron e-mails. Here is a map of a week's e-mail patterns in May 2001, when a new name suddenly appeared. Scientists found that this week's pattern differed greatly from others, suggesting

## The Alias

- k..allen did not appear in any prior graph.
- Perusal of the content of the emails determines that these were sent by Phillip Allen.
- phillip.allen appears in the previous graphs.
- A matched filter comparing neighborhoods was implemented and it found the correct match.
- In this work, we develop a "social space" version of the matched filter.


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## Aliases

- Given two graphs $G_{t}$ and $G_{t+1}$.
- Suppose we know some of the vertices are shared by these graphs (and which ones they are).
- There is one vertex in $G_{t+1}$ that we have not seen before.
- Assuming that this vertex appeared in $G_{t}$ with a different label, can we determine this vertex?


## Aliases

- Setup:
- Two graphs, $G_{t}=\left(V \cup U_{t}, E_{t}\right)$ and $G_{t+1}=\left(V \cup U_{t+1}, E_{t+1}\right)$.
- All vertices are labeled (email addresses).
- Vertices in $V$ are named (individual associated with the address).
- Vertices in $U_{i}$ are not named.
- Want to associate the names to the vertices in $U_{t+1}$.


## Methodology

- Assign the name to vertex $u$ whose vector $x_{v}$ is closest to the vector $x_{u}$.
- Optimize:

$$
\begin{aligned}
\left(X, Y_{1}, Y_{2}\right)=\arg \min _{X, Y_{1}, Y_{2}} & \left\|\left(\binom{X}{Y_{1}}\binom{X}{Y_{1}}^{T}\right)_{0}-A_{1}\right\|_{F}+ \\
& \left\|\left(\binom{X}{Y_{2}}\binom{X}{Y_{2}}^{T}\right)_{0}-A_{2}\right\|_{F}
\end{aligned}
$$

- $M_{O}$ means $M$ with the diagonal replaced with zeros.
- Thus, we are attempting to fit a set of vectors to the known and a set each for the unknown in the two graphs. Fitting to the knowns constrains the $Y_{i}$ to lie in the same space.


## The Setup

- Input $A_{1}, A_{2}$, the adjacency matrices of the graphs corresponding to the vertices $\left(V, U_{i}\right)$.
- Set $B$ to be the average of $A_{1}[V]$ and $A_{2}[V]$, the blocks corresponding to $V$.
- Set $N=n+n_{1}+n_{2}$.
- Set $A$ to be the $N \times N$ matrix with first $n \times n$ block equal to $B$, and blocks $A\left[V, U_{i}\right]=A_{i}, A\left[U_{i}, V\right]=A_{i}^{\prime}$.

$$
A=\left(\begin{array}{ccc}
\frac{A_{1}[V, V]+A_{2}[V, V]}{2} & A_{1}\left[V, U_{1}\right] & A_{2}\left[V, U_{2}\right] \\
A_{1}\left[U_{1}, V\right] & A_{1}\left[U_{1}, U_{1}\right] & Y^{\prime} \\
& Y & A_{2}\left[U_{2}, U_{2}\right]
\end{array}\right)
$$

where $Y$ is the dot product of vectors derived from $U_{1}$ and $U_{2}$.

## Fitting the Alias

1. Setup as described previously.
2. Set $D=0_{N \times N}$.
3. Set the first $n \times n$ block of $D$ equal to the the dot product of the result of running the least squares Algorithm on $B$.
3.1 While(Not Converged)
$3.2 Y=g_{d}(A+D)$
3.3 Set the unknown entries of $D$ (such as those corresponding to $U_{1} \times U_{2}$ ) to the dot products of the appropriate parts of $Y$.
4. Output $Y$

- Use the vectors to find the alias: closest named vector to the one associated with the alias.


## Alias Identification: k..allen $\rightarrow$ phillip.allen



Enron Executive

## Cartoon



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## Conclusions

- Social space provides a mechanism for modeling and inference on graphs and time series of graphs.
- Dot product graph model is simple, but easy to fit using linear algebra.
- Sparse matrix approaches can make this efficient:
- There appears to be an $O\left(n^{s}\right), 2<s<3$ matrix multiply in the algorithm, in order to determine the stopping criterion (compute the error).
- Some tricks can be played to reduce this for this application.
- By using only the change in the diagonal for determining convergence, we eliminate the need for the full matrix multiply, replacing it with an $O(n)$ operation. Note that we only need to check the diagonal, since once this stops changing the algorithm produces a fixed point.
- It is possible to add covariates (measurements at the nodes) into the model and still use the linear algebra approach, but this work is preliminary.


## Questions?

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